

A trifle on the stability of $f.x \leq g.y$
 (for the record)

We consider relation

$$R: f.x \leq g.y ,$$

in which f is an integer function not depending on y and g is an integer function not depending on x . We ask ourselves under what conditions statement

$$S: y := y + Y$$

maintains R . To that end we observe

$$\begin{aligned} & \text{wlp. } S.R \\ = & \quad \{ \text{axiom of assignment} \} \\ & f.x \leq g.(y + Y) \\ \Leftarrow & \quad \{ \text{using that } R \text{ is a precondition} \} \\ & g.y \leq g.(y + Y) \end{aligned}$$

- $\Leftarrow \quad \{ \text{assuming monotonicity of } g \}$
- $y \leq y + Y$
- $= \quad \{ \text{calc.} \}$
- $0 \leq Y$

- $\Leftarrow \quad \{ \text{assuming antimonicity of } g \}$
- $y + Y \leq y$
- $= \quad \{ \text{calc.} \}$
- $Y \leq 0$

As a result, increments of y maintain R for monotonic g and decrements of y maintain R for antimonic g . For changes in x we have similar requirements for f .

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In the development of multiprograms relations of the form $f.x \leq g.y$ are very wanted because of the above stability properties and -hence- abound. I think that in an orderly treatment of multiprogramming the above theorem has to be mentioned explicitly, in spite of its very simplicity. Hence this note.

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